# Detecting global optimality and extracting solutions in GloptiPoly

Didier HENRION<sup>1,2</sup> Jean-Bernard LASSERRE<sup>1</sup>

<sup>1</sup>LAAS-CNRS Toulouse <sup>2</sup>ÚTIA-AVČR Prague



Part 1 Description of GloptiPoly

# Brief description

GloptiPoly is written as an open-source, general purpose and user-friendly Matlab software

Optionally, problem definition made easier with Matlab Symbolic Math Toolbox, gateway to Maple kernel

Gloptipoly solves small to medium non-convex global optimization problems with multivariable real-valued polynomial objective functions and constraints

Software and documentation available at

www.laas.fr/~henrion/software/gloptipoly

# Metholodogy

GloptiPoly builds and solves a hierarchy of successive convex linear matrix inequality (LMI) relaxations of increasing size, whose optima are guaranteed to converge asymptotically to the global optimum



Relaxations are built from LMI formulation of sum of squares (SOS) decomposition of positive multivariable polynomials

LMIs solved with Jos Sturm's SeDuMi

In practice convergence is ensured fast, typically at 2nd or 3rd LMI relaxation

## LMI relaxation technique

Polynomial optimization problem

min 
$$g_0(x)$$
  
s.t.  $g_k(x) \ge 0, \quad k = 1, ..., m$ 

When  $p^*$  is the global optimum, SOS representation of positive polynomial

$$g_0(x) - p^* = q_0(x) + \sum_{k=1}^m g_k(x)q_k(x) \ge 0$$

where unknowns  $q_k(x)$  are SOS polynomials similar to Karush/Kuhn/Tucker multipliers

Using LMI representation of SOS polynomials successive LMI relaxations are obtained by increasing degrees of sought polynomials  $q_k(x)$ 

Theoretical proof of asymptotic convergence... ..but no tight degree upper bounds (not yet)

## LMI relaxations: illustration

Non-convex quadratic problem

$$\begin{array}{ll} \max & 2x_1^2 + 2x_2^2 - 2x_1x_2 - 2x_1 - 6x_2 + 10 \\ \text{s.t.} & -x_1^2 + 2x_1 \ge 0 \\ & -x_1^2 - x_2^2 + 2x_1x_2 + 1 \ge 0 \\ & -x_2^2 + 6x_2 - 8 \ge 0. \end{array}$$

LMI relaxation built by replacing each monomial  $x_1^i x_2^j$  with a new decision variable  $y_{ij}$ 

For example, quadratic expression

$$-x_1^2 - x_2^2 + 2x_1x_2 + 1 \ge 0$$

replaced with linear expression

$$-y_{20} - y_{02} + 2y_{11} + 1 \ge 0$$

New decision variables  $y_{ij}$  satisfy non-convex relations such as  $y_{10}y_{01} = y_{11}$  or  $y_{20} = y_{10}^2$ 

# LMI relaxations: illustration (2)

Relax these non-convex relations by enforcing LMI constraint

$$M_1^1(y) = \begin{bmatrix} 1 & y_{10} & y_{01} \\ y_{10} & y_{20} & y_{11} \\ y_{01} & y_{11} & y_{02} \end{bmatrix} \ge 0$$

Moment or measure matrix of first order relaxing monomials of degree up to 2

We remove the rank constraint on matrix  $M_1^1(y)$ 

First LMI relaxation of original global optimization problem is given by

$$\begin{array}{ll} \max & 2y_{20} + 2y_{02} - 2y_{11} - 2y_{10} - 6y_{01} + 10 \\ \text{s.t.} & -y_{20} + 2y_{10} \geq 0 \\ & -y_{20} - y_{02} + 2y_{11} + 1 \geq 0 \\ & -y_{02} + 6y_{01} - 8 \geq 0 \\ & M_1^1(y) \geq 0 \end{array}$$

## LMI relaxations: illustration (3)

To build second LMI relaxation, we must increase size of moment matrix so that it captures expressions of degrees up to 4

Second order moment matrix reads

	$\begin{bmatrix} 1 \end{bmatrix}$	$y_{10}$	$y_{01}$	$y_{20}$	$y_{11}$	$y_{02}$
$M_2^2(y) =$	$y_{10}$	$y_{20}$	$y_{11}$	$y_{30}$	$y_{21}$	$y_{12}$
	$y_{01}$	$y_{11}$	$y_{ m 02}$	$y_{21}$	$y_{12}$	$y_{OS}$
	$y_{20}$	$y_{30}$	$y_{21}$	$y_{ m 40}$	$y_{31}$	$y_{22}$
	$y_{11}$	$y_{21}$	$y_{12}$	$y_{31}$	$y_{22}$	$y_{13}$
	$y_{02}$	$y_{12}$	$y_{03}$	$y_{22}$	$y_{13}$	$y_{04}$ _

Constraints are also relaxed with additional variables Second LMI features feasible set included in first LMI feasible set, thus providing a tighter relaxation

# Numerical example (1)

Quadratic problem 3.5 in [Floudas/Pardalos 99]

$$\begin{array}{rl} \min & -2x_1 + x_2 - x_3 \\ \text{s.t.} & x_1(4x_1 - 4x_2 + 4x_3 - 20) + x_2(2x_2 - 2x_3 + 9) \\ & + x_3(2x_3 - 13) + 24 \ge 0 \\ & x_1 + x_2 + x_3 \le 4, \quad 3x_2 + x_3 \le 6 \\ & 0 \le x_1 \le 2, \quad 0 \le x_2, \quad 0 \le x_3 \le 3. \end{array}$$

To define this problem with GloptiPoly we use the following Matlab/Maple script

To solve the first LMI relaxation we type

```
>> output = gloptipoly(P)
output =
    status: 0
    crit: -6.0000
    sol: {}
```

Field status = 0 indicates that it is not possible to detect global optimality with this LMI relaxation, hence crit = -6.0000 is a lower bound on the global optimum

## Numerical example (2)

Next we try to solve the second, third and fourth LMI relaxations

```
>> output = gloptipoly(P,2)
                                     >> output = gloptipoly(P,3)
output =
                                     output =
    status: 0
                                         status: 0
      crit: -5.6923
                                           crit: -4.0685
                                            sol: {}
       sol: {}
>> output = gloptipoly(P,4)
output =
    status: 1
      crit: -4.0000
       sol: {[3x1 double] [3x1 double]}
>> output.sol{:}
ans =
                                     ans =
    2.0000
                                         0.5000
    0.0000
                                         0.0000
    0.0000
                                         3.0000
```

Both second and third LMI relaxations return tighter lower bounds on the global optimum

Eventually global optimality is reached at fourth LMI relaxation (certified by status = 1)

GloptiPoly also returns two globally optimal solutions:

$$x_1 = 2, x_2 = 0, x_3 = 0$$

and

$$x_1 = 0.5, x_2 = 0, x_3 = 3$$

leading to

crit = -4.0000

# Numerical example (3)

Number of LMI variables (M) and size of relaxed LMI problem (N) increase quickly with relaxation order:

Relaxation	LMI opt	M	N
1	-6.0000	9	24
2	-5.6923	34	228
3	-4.0685	83	1200
4	-4.0000	164	4425
5	-4.0000	285	12936
6	-4.0000	454	32144

..yet fourth LMI relaxation was solved in about 2.5 seconds on a PC Pentium IV 1.6 MHz

## Complexity

d: overall polynomial degree  $(2\delta = d \text{ or } d + 1)$ m: number of polynomial constraints n: number of polynomial variables M: number of LMI decision variables N: size of LMI

$$M = \begin{pmatrix} n+2\delta \\ 2\delta \end{pmatrix} - 1$$
$$N = \begin{pmatrix} n+\delta \\ \delta \end{pmatrix} + m \begin{pmatrix} n+\delta-1 \\ \delta-1 \end{pmatrix}$$

When n is fixed:

- M grows polynomially in  $O(\delta^n)$
- N grows polynomially in  $O(m\delta^n)$

# Features

General features of GloptiPoly:

- Certificate of global optimality (rank checks)
- Automatic extraction of globally optimal solutions (multiple eigenvectors)
- 0-1 or  $\pm 1$  integer constraints on some of the decision variables (combinatorial optimization problems)
- Generation of input and output data in SeDuMi's format
- Generation of moment matrices associated with LMI relaxations (rank checks)
- User-defined scaling of decision variables (to improve numerical behavior)
- Exploits sparsity of polynomial data

# Benchmark examples Continuous problems

Mostly from Floudas/Pardalos 1999 handbook

About 80 % of pbs solved with LMI relaxation of small order (typically 2 or 3) in less than 3 seconds on a PC Pentium IV at 1.6 MHz with 512 Mb RAM



Benchmark exmaples Discrete problems

From Floudas/Pardalos handbook and also Anjos' Ph.D (Univ Waterloo)

By perturbing criterion (destroys symmetry) global convergence ensured on 80 % of pbs in less than 4 seconds



MAXCUT on antiweb  $AW_9^2$  graph

# Benchmark examples Polynomial systems of equations

From Verschelde's database and Frisco INRIA project Real coefficients & solutions only

Out of 59 systems:

- 61 % solved in t < 10 secs
- 20 % solved in 10 < t < 100 secs
- 10 % solved in t  $\geq$  100 secs
- 9 % out of memory

No criterion optimized No enumeration of all solutions



Intersections of seventh and eighth degree polynomial curves

# Part 2 Extracting solutions

# Detecting global optimality

Global optimization problem

$$p^{\star} = \min_{x} g_{0}(x)$$
  
s.t.  $g_{i}(x) \ge 0, i = 1, 2..$ 

Let deg  $g_i(x) = 2d_i - 1$  or  $2d_i$  and  $d = \max_i d_i$ 

LMI relaxation of order  $\boldsymbol{k}$ 

$$p_k^* = \min_y \sum_lpha (g_0)_lpha y_lpha \ ext{ s.t. } M_k(y) \succeq 0 \ M_{k-d_i}(g_iy) \succeq 0, \ i=1,2..$$

with solution  $y_k^*$ 

Hierarchy with convergence guarantee

 $p_d^* \le p_{d+1}^* \le \dots \le p_{k^*}^* = p^*.$ 

for (generally) small  $k^*$ 

Sufficient condition for global optimality rank of moment matrices

$$\operatorname{rank} M_k(y_k^*) = \operatorname{rank} M_{k-d}(y_k^*)$$

Need for extraction procedures

Because rank condition is only sufficient any other global optimality certificate is welcome

Solution extraction procedure suggested by Arnold Neumaier

Inspired by Chesi, Garulli, Tesi, Vicino (IEEE CDC 2000)

Based on the work by Corless, Gianni, Trager (ACM ISSAC 1997)

Key ideas:

- Cholesky decomposition of moment matrix
- column reduced echelon form
- computation of common eigenvalues

# Who is Cholesky ?

André Louis Cholesky (1875-1918) was a French military officer (graduated from Ecole Polytechnique) involved in geodesy

He proposed a new procedure for solving least-squares triangulation problems

He fell for his country during World War I



Work posthumously published in

Commandant Benoît. Procédé du Commandant Cholesky. *Bulletin Géodésique*, No. 2, pp. 67-77, Toulouse, Privat, 1924.

#### Nice biography in

C. Brezinski. André Louis Cholesky. *Bulletin of the Belgian Mathematical Society*, Vol. 3, pp. 45-50, 1996.

#### I. - NOTICE5 SCIENTIFIQUES

Commandant BENOIT'.

NOTE SUR UNE MÉTHODE DE RÉSOLUTION DES ÉQUA-TIONS NORMALES PROVENANT DE L'APPLICATION DE LA MÉTHODE DES MOINDRES CARRÉS A UN SYSTÈME D'ÉQUATIONS LINÉAIRES EN NOMBRE INFÉRIEUR A CELUI DES INCONNUES. — APPLICATION DE LA MÉ-THODE A LA RÉSOLUTION D'UN SYSTÈME DEFINI D'ÉQUATIONS LINÉAIRES.

(Procédé du Commandant Cholesky !.)

Le Commandant d'Artillerie Cholesky, du Service géographique de l'Armée, tué pendant la grande guerre, a imaginé, au cours de recherches sur la compensation des réseaux géodésiques, un procédé très ingénieux de résolution des équations dites normales, obtenues par application de la méthode des moindres carrés à des équations linéaires en nombre inférieur à celui des inconnues. Il en a conclu une méthode générale de résolution des équations linéaires.

Nous suivrons, pour la démonstration de cette méthode, la progression môme qui a servi au Commandant Cholesky pour l'imaginer.

1. De l'Artillèrie coloniale, ancien officier géodésien au Service géographique de l'Armée et au Service géographique de l'Indo-Chine, Membre du Comité national français de Géodésie et Géophysique.

2. Sur le Commandant Cholesky, tué à l'ennemi le 31 août 1918, voir la notice biographique insérée dans le volume du Balletin géodésique de 1922 intitulé : Union géodésique et géophysique internationale, Première Assemblée générale, Rome, mai 1992, Section de Géodésie, Toulouse, Privat, 1922, in-8°, 241 p., pp., 159 à 161.

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# Cholesky factorization

Extract Cholesky factor V of moment matrix

 $M_k(y_k^\star) = VV'$ 

Matrix V has r columns, corresponding to r globally optimal solutions  $x_j^*$ , j = 1, 2, ..., r (provided global optimum was reached)

Denote

 $v = \begin{bmatrix} 1 & x_1 & x_2 & \dots & x_n & x_1^2 & x_1 & x_2 & \dots & x_1 & x_n & x_2^2 & x_2 & x_3 & \dots & x_n^k & \dots & x_n^k \end{bmatrix}'$ a basis for polynomials of degree at most k

By definition of the moment matrix:

$$M_k(y_k^\star) = V^\star(V^\star)'$$

where

$$V^{\star} = \left[\begin{array}{cccc} v_1^{\star} & v_2^{\star} & \cdots & v_r^{\star} \end{array}\right]$$

and  $v_i^{\star}$  is polynomial basis v evaluated at solution  $x_i^{\star}$ 

Extracting solutions amounts to finding linear transformation between Cholesky factors V and  $V^*$ 

#### Reduction to column echelon form

Next step is reduction of V into column echelon form

$$U = \begin{bmatrix} 1 & & & \\ * & & & \\ 0 & 1 & & \\ 0 & 0 & 1 & & \\ * & * & * & \\ \vdots & & \ddots & \\ 0 & 0 & 0 & \cdots & 1 \\ * & * & * & \cdots & * \\ \vdots & & & \vdots \\ * & * & * & \cdots & * \end{bmatrix}$$

by Gaussian elimination with column pivoting

Each row in U = monomial  $x_{\alpha}$  in basis v

Pivot entry in U = monomial  $x_{\beta_j}$  in generating basis of the set of solutions

In other words, denoting

$$w = \begin{bmatrix} x_{eta_1} & x_{eta_2} & \dots & x_{eta_r} \end{bmatrix}'$$

it holds

v = Uw

for all solutions  $x_j^*$ ,  $j = 1, 2, \ldots, r$ 

# Multiplication matrices

For each first degree monomial  $x_i$  extract from U the r-by-r multiplication matrix  $N_i$  containing coefficients of product monomials  $x_i x_{\beta_i}$  in generating basis w, i.e. such that

 $N_i w = x_i w \quad i = 1, 2, \dots, n$ 

If monomial  $x_i x_{\beta_j}$  is not represented in U, then extraction algorithm fails and order k of LMI relaxation must be increased

Given matrices  $N_i$  finding scalars  $x_i$  is an eigenvalue problem

Extracting solutions amounts to solving an eigenvalue problem

Eigenvectors w are shared by matrices  $N_i$  so it is a particular common eigenvalue problem

# Common eigenvalue problem

Build combination of multiplication matrices

$$N = \sum_{i=1}^{n} \lambda_i N_i$$

where  $\lambda_i$  are random positive numbers (summing up to one)

Compute ordered Schur decomposition

$$N = QTQ'$$

where

$$Q = \left[\begin{array}{cccc} q_1 & q_2 & \cdots & q_r \end{array}\right]$$

is orthogonal and  ${\boldsymbol{T}}$  upper triangular

Finally, due to orthogonality of vectors  $q_i$ , *i*th entry in solution vector  $x_i^{\star}$  is given by

$$(x_j^*)_i = q_j' N_i q_j, \ i = 1, 2, \dots, n, \ j = 1, 2, \dots, r$$

# Number of solutions

No easy way to control number of extracted solutions in case of multiple global optima

Number of solutions = rank of moment matrix, but enforcing rank in an LMI is a difficult nonconvex problem

By default GloptiPoly minimizes the trace (sum of eigenvalues) of the moment matrix, which may indirectly minimize the rank (number of non-zero eigenvalues)

Practical experiments reveal that low rank moment matrices ensure faster convergence of LMI relaxations to global optimum

#### First example

Non-convex quadratic optimization

$$p^* = \max_x (x_1 - 1)^2 + (x_1 - x_2)^2 + (x_2 - 3)^2$$
  
s.t.  $(x_1 - 1)^2 \le 1$   
 $(x_1 - x_2)^2 \le 1$   
 $(x_2 - 3)^2 \le 1$ 

First LMI relaxation yields  $p_1^{\star} = -3$  and rank  $M_1(y^{\star}) = 3$ , extraction algorithm fails due to incomplete monomial basis

Second LMI relaxation yields  $p_2^{\star} = -2$  and

rank  $M_1(y^*) = \operatorname{rank} M_2(y^*) = 3$ 

so rank condition ensures global optimality

#### First example: Cholesky factor

# Moment matrix of order k = 2 reads

, 0.200. 0.02.0 12.000 10.000 12.1000 10.1000 1	$M_2(y^*) =$	1.0000 1.5868 2.2477 2.7603 3.6690 5.2387	$\begin{array}{c} 1.5868 \\ 2.7603 \\ 3.6690 \\ 5.1073 \\ 6.5115 \\ 8.8245 \end{array}$	2.2477 3.6690 5.2387 6.5115 8.8245 12.7072	2.7603 5.1073 6.5115 9.8013 12.1965 15.9960	3.6690 6.5115 8.8245 12.1965 15.9960 22.1084	5.2387 8.8245 12.7072 15.9960 22.1084 32.1036
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## Positive semidefinite with rank 3

## Cholesky factor

$$V = \begin{bmatrix} -0.9384 & -0.0247 & 0.3447 \\ -1.6188 & 0.3036 & 0.2182 \\ -2.2486 & -0.1822 & 0.3864 \\ -2.9796 & 0.9603 & -0.0348 \\ -3.9813 & 0.3417 & -0.1697 \\ -5.6128 & -0.7627 & -0.1365 \end{bmatrix}$$

satisfies

$$M_k(y_k^\star) = VV'$$

#### First example: column echelon form

Gaussian elimination on V yields

$$U = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ -2 & 3 & 0 \\ -4 & 2 & 2 \\ -6 & 0 & 5 \end{bmatrix} \begin{bmatrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1 x_2 \\ x_2^2 \\ 1 \end{bmatrix}$$

which means that solutions to be extracted satisfy the system of polynomial equations

$$\begin{array}{rcrr} x_1^2 &=& -2 + 3x_1 \\ x_1x_2 &=& -4 + 2x_1 + 2x_2 \\ x_2^2 &=& -6 + 5x_2 \end{array}$$

in polynomial basis 1,  $x_1$ ,  $x_2$ 

#### First example: extraction

Multiplication matrices of monomials  $x_1$  and  $x_2$  in polynomial basis 1,  $x_1$ ,  $x_2$  are extracted from U:

$$N_1 = \begin{bmatrix} 0 & 1 & 0 \\ -2 & 3 & 0 \\ -4 & 2 & 2 \end{bmatrix}, \quad N_2 = \begin{bmatrix} 0 & 0 & 1 \\ -4 & 2 & 2 \\ -6 & 0 & 5 \end{bmatrix}$$

Random linear combination

$$N = 0.6909N_1 + 0.3091N_2$$

Schur decomposition of N = QTQ' yields

$$Q = \begin{bmatrix} 0.4082 & 0.1826 & -0.8944 \\ 0.4082 & -0.9129 & -0.0000 \\ 0.8165 & 0.3651 & 0.4472 \end{bmatrix}$$

Projections of orthogonal columns of Q onto N yield the 3 expected globally optimal solutions

$$x_1^* = \begin{bmatrix} 1\\2 \end{bmatrix} x_2^* = \begin{bmatrix} 2\\2 \end{bmatrix} x_3^* = \begin{bmatrix} 2\\3 \end{bmatrix}$$

## Second example: minimum trace LMI

Polynomial system of equations

$$\begin{array}{rcl}
x_1^2 + x_2^2 &=& 1\\ x_1^3 + (2 + x_3)x_1x_2 + x_2^3 &=& 1\\ & & x_3^2 &=& 2 \end{array}$$

No objective function, so GloptiPoly minimizes the trace of the moment matrix

Extraction on 2nd LMI relaxation fails due to incomplete basis

3rd LMI relaxation yields two globally optimal solutions

$$x_1^* = \begin{bmatrix} 0.5826\\ -0.8128\\ -1.4142 \end{bmatrix} \quad x_2^* = \begin{bmatrix} -0.8128\\ 0.5826\\ -1.4142 \end{bmatrix}$$

Second example: zero objective function

With zero objective function GloptiPoly at the 3rd LMI relaxation yields

rank $M_1(y^*) = 4 \neq \text{rank}M_2(y^*) = \text{rank}M_3(y^*) = 6$ so rank condition cannot ensure optimality

However, extraction algorithm returns 6 globally optimum solutions

$$x_{1}^{*} = \begin{bmatrix} -0.8128\\ 0.5826\\ -1.4142 \end{bmatrix} \quad x_{2}^{*} = \begin{bmatrix} 0.5826\\ -0.8128\\ -1.4142 \end{bmatrix}$$
$$x_{3}^{*} = \begin{bmatrix} 0.0000\\ 1.0000\\ -1.4142 \end{bmatrix} \quad x_{4}^{*} = \begin{bmatrix} 1.0000\\ 0.0000\\ -1.4142 \end{bmatrix}$$
$$x_{5}^{*} = \begin{bmatrix} 0.0000\\ 1.0000\\ 1.4142 \end{bmatrix} \quad x_{6}^{*} = \begin{bmatrix} 1.0000\\ 0.0000\\ 1.4142 \end{bmatrix}$$

thus proving global optimality of LMI

#### Motzkin-like polynomial

Polynomial

$$\frac{1}{27} + x^2 y^2 (x^2 + y^2 - 1)$$

vanishes at  $|x| = |y| = \sqrt{3}/3$  and remains globally non-negative for real x and ybut cannot be written as an SOS



# Some kind of magic in GloptiPoly

GloptiPoly finds approximate SOS decomposition of Motzkin polynomial

With 8th LMI relaxation we obtain

$$\frac{1}{27} + x^2 y^2 (x^2 + y^2 - 1) = \sum_{i=1}^{32} a_i^2 q_i^2 (x, y) + \varepsilon r(x, y)$$

where 
$$\|q_i\|_2 = \|r\|_2 = 1$$
  
and  $\varepsilon \leq 10^{-8} < a_i^2$ , deg  $q_i \leq 8$ 

Cone of SOS polynomials is dense in set of polynomials nonnegative over box [-1,1]

Numerical inaccuracy helps finding higher degree SOS polynomial close to Motzkin polynomial



Constrained Motzkin polynomial

Additional redundant constraint

$$x^2 + y^2 \le R^2$$

with  $R^2 > 2/3$  (includes the 4 global minima)

For R = 1 at the 3rd LMI relaxation we obtain

$$\frac{1}{27} + x^2 y^2 (x^2 + y^2 - 1) = \sum_{i=1}^{6} a_i^2 q_i^2 (x, y) + (R^2 - x^2 - y^2) \sum_{i=1}^{2} b_i^2 r_i^2 (x, y)$$

where deg  $q_i \leq 3$ , deg  $r_i \leq 2$ 

$$R^2$$
 1
 2
 3
 4
  $\cdots$ 
 $\infty$ 

 LMI
 3
 4
 5
 6
  $\cdots$ 
 8

Relevance of feasibility radius in SDP solver and GloptiPoly

# Conclusions

GloptiPoly is a general-purpose software with a user-friendly interface

Pedagogical flavor, black-box approach, no expert tuning required to cope with very distinct applied maths and engineering pbs

Automatic detection of global optimality and extraction of solutions

Not a competitor to highly specialized codes for solving polynomial systems of equations or large combinatorial optimization pbs

Numerical conditioning (Chebyshev basis) and problem structure (Hankel/Toeplitz matrices) deserve further study

See also Parrilo's **SOSTOOLS** software

# Further news

# Major extension of GloptiPoly planned (hopefully) for winter 2003

- Matlab classes for multivariate polynomials and moment matrices
- general moment problems
- performance analysis for stochastic systems in ecology and finance
- robust control problems
- relaxations of robust LMIs

# Research efforts

- bilinearity in decision variables
- tailored interior-point algorithms

Regularly updated information at

www.laas.fr/~henrion www.laas.fr/~lasserre